

Quiz 2 solution

Cmps 211

Question 2

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(a) problem definition

input $\{a_j\}_{j=1}^n$ ^{size of input} $'<'$ _{order is defined for comparison}

Output let $P(i) = "a_i < \prod_{k=1}^{i-1} a_k"$

\therefore return $A = \{a_i \mid P(i), 1 \leq i \leq n\}$

(b) paradigm + pseudo-code

Paradigm = iterative

Pseudo code =

lessThanProd($\{a_j\}_{j=1}^n, '<'$)

{ $A \leftarrow \emptyset$;

$prod \leftarrow 1$;

 for ($j = 2$; $j \leq n$; $j++$)

 { $prod \leftarrow prod * a_{j-1}$;

 if $a_j < prod$

$A \leftarrow A \cup \{a_j\}$

 }
 return A
}

c) Best/worst case definition

Input size : n

* Best case : $A = \emptyset \quad \forall j \geq 2 \neg P(j)$

* worst case : $A = \{a_j\}_{j=2}^n : \forall j \geq 2 P(j)$

d) Space requirement

* $S(n) = \overset{\text{read input}}{n} + \overset{\text{set } A \text{ returns}}{n} + \overset{\text{for prod variable}}{1} + \overset{\text{every read/write instruction requires space}}{1}$
for j

$$= 2n + 2$$

$$= 2(n+1)$$

$$= O(n)$$

e) Worst case runtime

$T(n) = \underset{\text{assign } A}{1} + \underset{\text{assign prod}}{1} + \overset{\text{\# of iterations}}{(n-1)} \left(\underset{\text{\# cost per iteration}}{5} \right) + \underset{\text{return } A}{1}$

$$= 2 + 5n - 5 + 1$$

$$= 5n - 2$$

8) order of growth

$$S_{n-2} = O(n) \\ = \Theta(n)$$

9) prove (P)

$$S_{n-2} = \Theta(n)$$

Prove $\exists c_1, c_2, n_0 : c_1 n \leq S_{n-2} \leq c_2 n \quad \forall n \geq n_0$

$n_0 = 1$	$1n \leq S_{n-2} \leq 7n$
$c_1 = 1$	$\therefore S_{n-2} = \Theta(n)$
$c_2 = 7$	

10) loop invariants

after the end of j^{th} iteration set A contains those elements \leq product of previous elements seen so far

(After the j^{th} iteration)

$$A(j) = A = \{ a_i \mid P(i) : 1 \leq i < j \}$$

index of iteration

(i) Proof by Induction

by Induction

Base cases ~~Recursion~~ assert $Q(1)$

$Q(1)$ - at end of 1st iteration $A = \{a_j \mid P(i) : 1 \leq i < 2\}$

before 1st loop we know

$$\text{prod} = 1$$

$$A = \phi$$

during if ~~prod~~ $a_2 < \text{prod} + a_2$

$$A = \{a_2\}$$

else

$$A = \phi$$

After

$A = \phi \vee A = \{a_2\} \therefore Q(1)$

Inductive Step

assert $Q(j)$ assuming $Q(j-1)$

before

assume $Q(j-1)$:

at j^{th} iteration $A = \{a \mid P(i) : 1 \leq i < j-1\}$

during

if $a_{j-1} < \text{prod} + a_{j-1}$

$$A = A \cup \{a_{j-1}\}$$

else

$$A = A$$

after

$A = A \vee A \cup \{a_{j-1}\} \therefore Q(j)$